Advanced C Programming Profiling

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Today

Profiling

Invasive Profiling Non-Invasive Profiling

Tools

gprof gcov valgrind oprofile

Conclusion

What is a Profiler?

Analyse the runtime behavior of the program

- ▶ Which parts (functions, statements, ...) of a program take how long?
- ▶ How often are functions called?
- Which functions call which
 - Construct the dynamic call graph
- Memory consumption
 - Memory accesses
 - memory leaks
 - Cache performance

Invasive Profiling

- Modify the program (code instrumentation)
- Insert calls to functions that record data

- Advantages:
 - Very precise
 - Theoretically at the instruction level
 - Precise call graph
- ► Disadvantages:
 - Potentially very high overhead
 - Depends on the instrumentation code that is inserted
 - Cannot profile already running systems (long running servers)
 - Can only profile application (not complete system)

Non-Invasive Profiling

- Statistic sampling of the program
- Use a fixed time interval or Hardware performance counters (CPU feature) to trigger sampling events
- Record instruction pointer at each sampling event
- Advantages:
 - Small overhead
 - Hardware assisted
 - Can profile the whole system (even the kernel!)
- Disadvantages:
 - not precise r "only" statistical data
 - Call Graph possibly not complete
 some functions are never sampled

Profiles

► Flat Profile

How much time does the program spend in which function?

Call Graph Which function calls which function how often?

► Annotated Sources
Annotate each source line with number of executions

gprof

Mixture of invasive and statistical profiling

Invasive Part

- ▶ gcc inserts calls to a function mcount into prologue of each function
- ▶ Compile with -g and -pg
- ▶ mcount can figure out its caller ☞ we can construct the call graph
- mcount counts the number of invocations for each function
- ► Call to mcount is the only instrumentation

 □ almost as efficient as normal build
- After program is run, there is a file called gmon.out containing profiling data
- ▶ Evaluate contents of gmon.out with gprof name-of-program

gprof

Statistical Part

- ► Kernel samples instruction pointer (IP) on each timer interrupt (100/s)
- Increments a counter in a histogram of address ranges
 cannot track the exact location where timer interrupt happened
- Provides a frequency distribution over code locations
- ▶ Beware of low samplerate
- Short running programs will mostly not provide meaningful data
- Accumulation of several profile runs is possible:

```
$ ./test_program
$ mv gmon.out gmon.sum
$ ./test_program
$ gprof -s ./test_program gmon.out gmon.sum
```

- Analyses coverage of program code
- Which line was executed how often
- Helps for finding code that
 - can profit from optimizations
 - that is not covered by test cases
- Use GCC flags
 - -fprofile-arcs: collect info about jumps
 - -ftest-coverage: collect info about code coverage
- Attention: Multiple code lines might be merged to one instruction

```
100: 12:if (a != b)

100: 13: c = 1;

100: 14:else

100: 15: c = 0;
```

valgrind

- ► JIT-compiler / translator:
 - Construct intermediate representation from x86 assembly code
 - Add instrumentation code
 - Compile back to x86
- ▶ Done while program is loaded
- Is not only a profiler!
- No compiler flags / recompilation needed (though -g -fno-inline advisable to analyse output)
- Program runtime can degrade drastically due to instrumentation code and recompilation
- can escape to debugger on certain events
 very handy when debugging memory leaks
- Disadvantage:
 - program might run an order of magnitude slower
 - program might consume an order of magnitude more memory

valgrind

Tools

memcheck

- ▶ Redirects calls to malloc and the like
- Keeps track of all allocated memory
- ▶ Instruments references to warn about "bad" memory accesses
 - uninitialized
 - already freed
- ► Detects memory leaks
- Warns about jumps taken upon uninitialized values

cachegrind

- ► Instruments memory accesses
- ▶ Simulates (!) a L1 and L2 cache in software
- ► Gives precise data about cache misses

callgrind

▶ Records the call graph

Hint

Use kcachegrind for visualization

oprofile

- ▶ Non-invasive
- Kernel module and user-space daemon
- Does not modify the program at all
- ▶ -g for debug symbols recommendable
- Sampling uses performance counters
- ... or timer interrupt of perf. counters not available
- Profiles the whole system (also the kernel!)
- Can distill data for each binary separately
- ► For Windows, use Intel vTune (\$\$\$)

oprofile

Performance Counter

- Set of hardware registers for a plethora of events
- ▶ Differ from processor model to another
- Very detailed events trackable. Examples:
 - ► L2 cache misses
 - Retired instructions
 - Outstanding bus requests
 - ... and many more
- Basic modus operandi:
 - Kernel module tells the CPU to fire an exception after a certain number of events of a certain type have occurred
 - CPU traps into kernel
 - instruction pointer is recorded in a buffer (no histograms)

oprofile

Howto

- ▶ Use opcontrol to control the daemon/module
- opcontrol --init to load module and daemon
- ▶ opcontrol -s to start sampling
- opcontrol -t to stop sampling
- ▶ opcontrol --dump flushes the event log
- ▶ opcontrol --list-events shows available performance counters
- opreport -1 prog-name gives breakdown of samples per function in prog-name

Conclusion

- Many different profiling methods exist
- ▶ gprof
 - is obsolete
 - use only to get a quick impression
 - and for the call graph
 - sampling might be too imprecise
- valgrind
 - easy to use
 - no recompile
 - precise
 - good visualization (kcachegrind)
 - but large increase in runtime
- ▶ oprofile
 - much more precise than gprof
 - can profile exotic machine events if you are going for the last cycles
 - not as precise as valgrind
 - need root rights on the machine